

Network Technologies (TCP/IP Suite)

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Outline

- ▲ Direct & Indirect delivery
- ▲ Forwarding & Forwarding Techniques
- ▲ Forwarding with classful addressing
 - Without subnetting
 - With subnetting
- ▲ Forwarding with classless addressing
 - Longest subnet mask
- ▲ Aggregation
- ▲ Autonomous Systems

Trivia – Network Elements?

▲ Repeater?

▲ Hub?

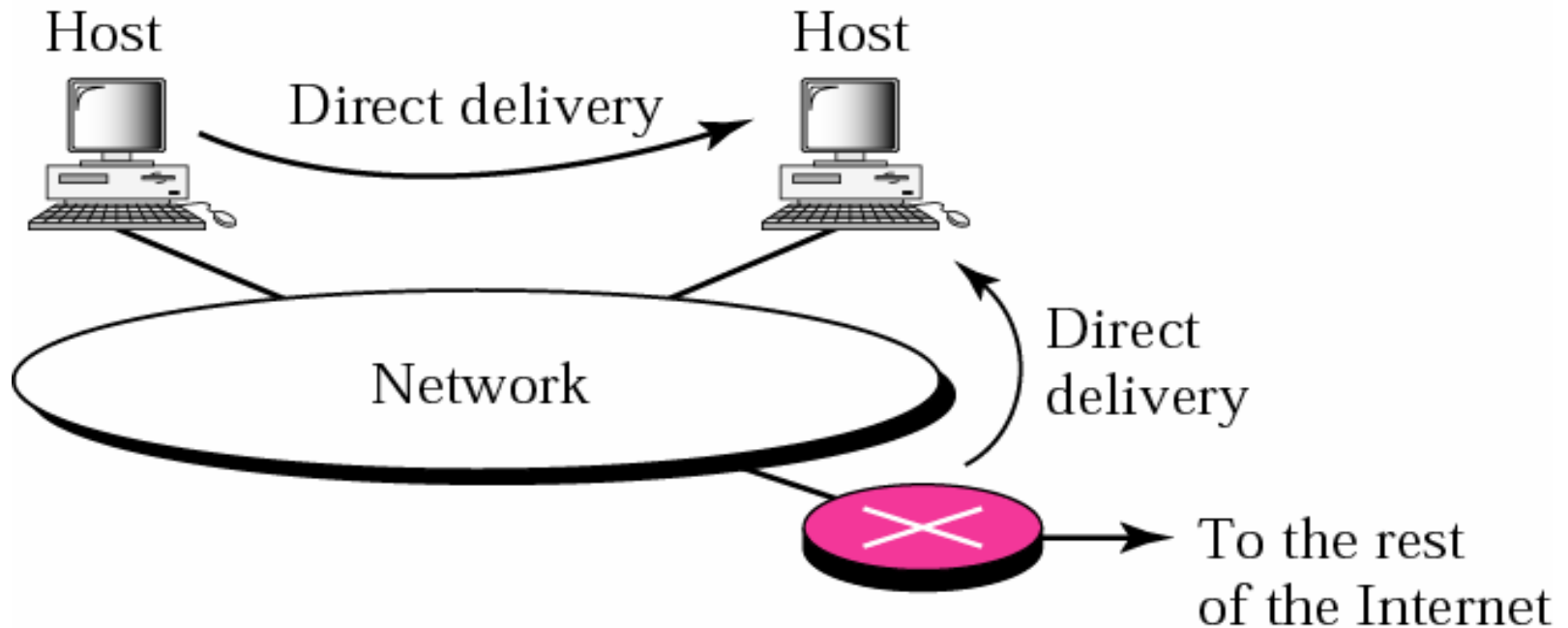
▲ Bridge?

▲ Switch?

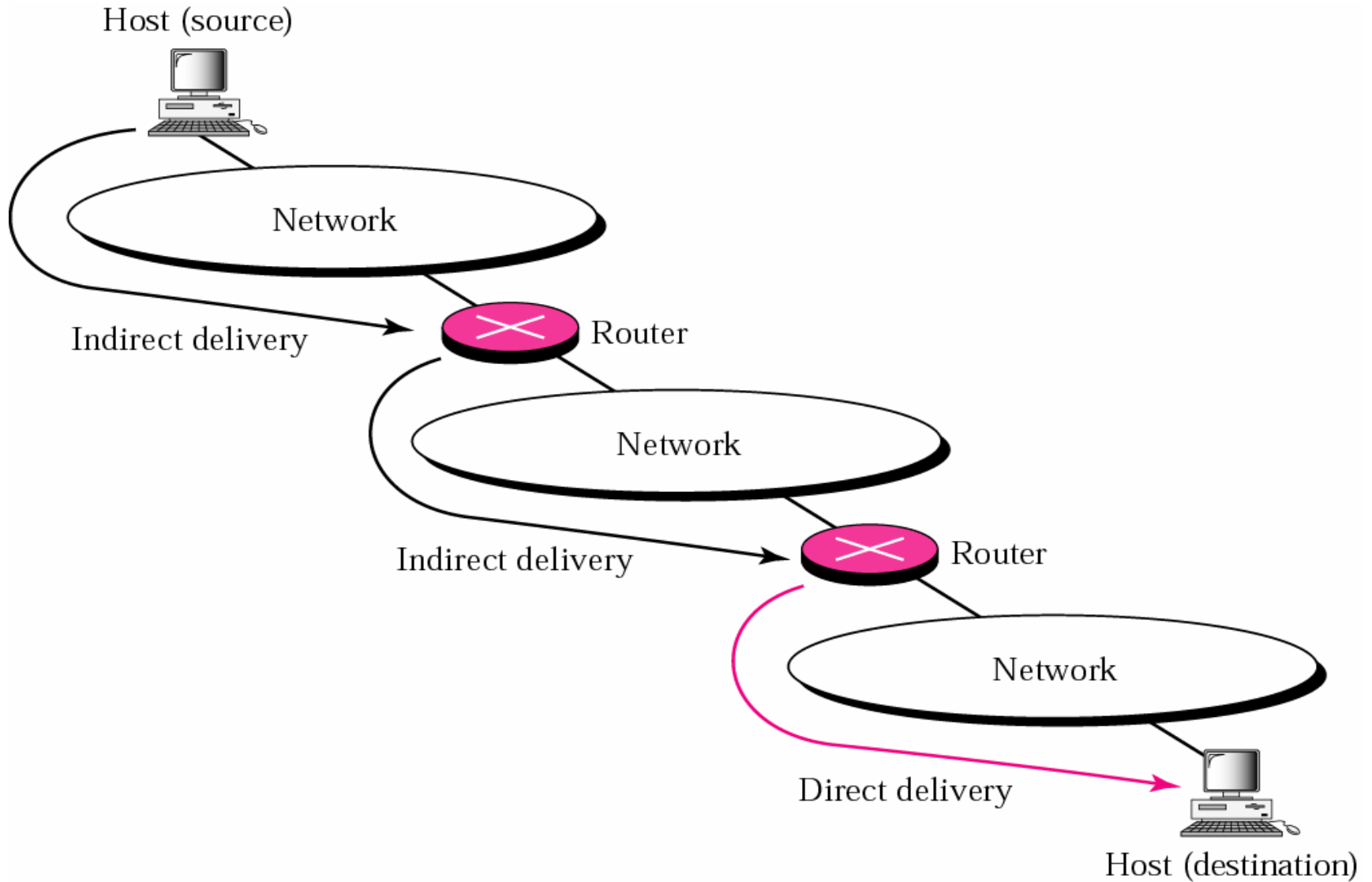
▲ Router?

– Layer-3-Switch?

Direct delivery



Indirect delivery



Forwarding VS Routing

- ▲ Earlier we assumed that routers have knowledge of network structure
 - Virtual Circuits ~ routing is an issue for the connection request
 - Datagram Model ~ routing is an issue for every packet
- ▲ Forwarding ~ collect a packet, look at the destination, consult a **table**, send the packet in the right direction
 - Relatively simple
- ▲ Routing ~ The process of constructing the **tables**
 - Reasonably complex distributed algorithms

Forwarding Tables VS Routing Tables

- ▲ Terms used interchangeably ~ incorrect!
- ▲ Forwarding tables ~ used when packet is being sent on the right path
 - Mapping of network identity to physical address (of the next hop on the path) and outgoing interface
- ▲ Routing tables ~ Built by the routing algorithms as a precursor to building the forwarding tables
 - Mapping of network identity to the next hop on the path

Forwarding Tables VS Routing Tables

▲ Whether these tables are separate data structures is a consideration for implementation

– Usually separate

| Network Identity | Next Hop |
|-------------------------|-----------------|
| 10 | 171.69.245.10 |

Routing table

| Network Identity | Interface | MAC Address |
|-------------------------|------------------|--------------------|
| 10 | if4 | 8:0:2b:e4:b:1:2 |

Forwarding table

Forwarding

- ▶ Forwarding means to place the packet in its route to its destination
- ▶ Forwarding requires a host or a router to have a routing table
- ▶ Whenever a packet is to be forwarded, we have to lookup in the routing table to identify the route to the destination

Forwarding techniques

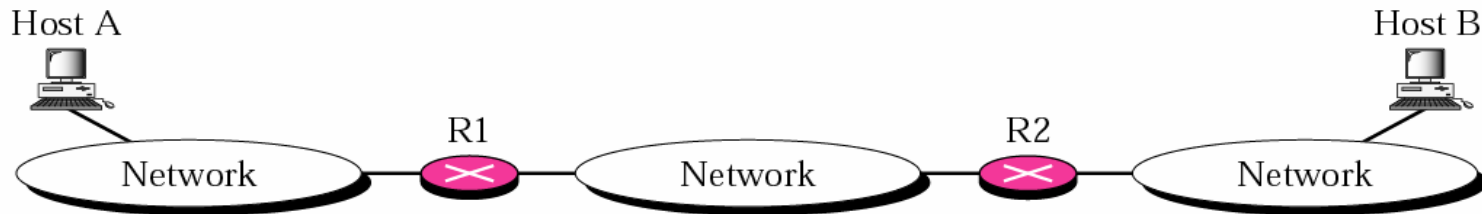
- ▶ Considering the number of interfaces and interconnecting devices, we need to have an efficient forwarding technique
 - We cannot maintain a list of all the hosts/interfaces
- ▶ Efficient algorithm
- ▶ A small routing table

Next-Hop Method

- ✦ The table holds only the address of the next hop
- ✦ No information about the complete route

| Routing table for host A | | Routing table for R1 | | Routing table for R2 | |
|--------------------------|----------------|----------------------|------------|----------------------|--------|
| Destination | Route | Destination | Route | Destination | Route |
| Host B | R1, R2, Host B | Host B | R2, Host B | Host B | Host B |

a. Routing tables based on route



| Routing table for host A | | Routing table for R1 | | Routing table for R2 | |
|--------------------------|----------|----------------------|----------|----------------------|----------|
| Destination | Next Hop | Destination | Next Hop | Destination | Next Hop |
| Host B | R1 | Host B | R2 | Host B | Ñ |

b. Routing tables based on next hop

Network-Specific Method

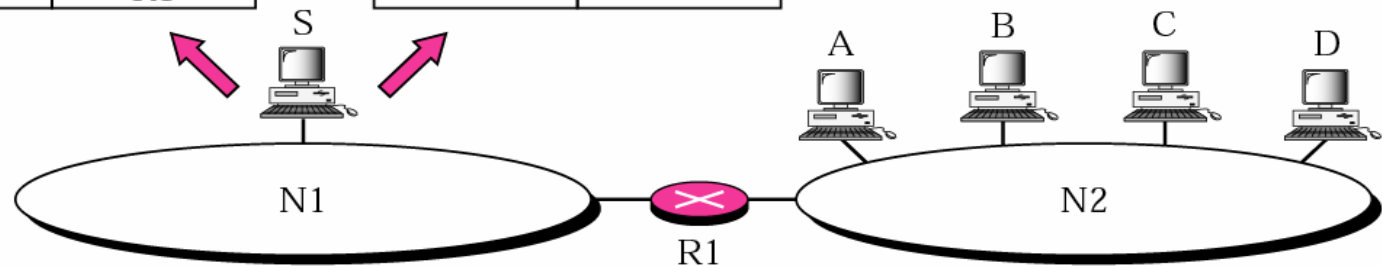
- For every destination host connected to the same network, we have only one entry that defines the address of the destination address

Routing table for host S based on host-specific method

| Destination | Next Hop |
|-------------|----------|
| A | R1 |
| B | R1 |
| C | R1 |
| D | R1 |

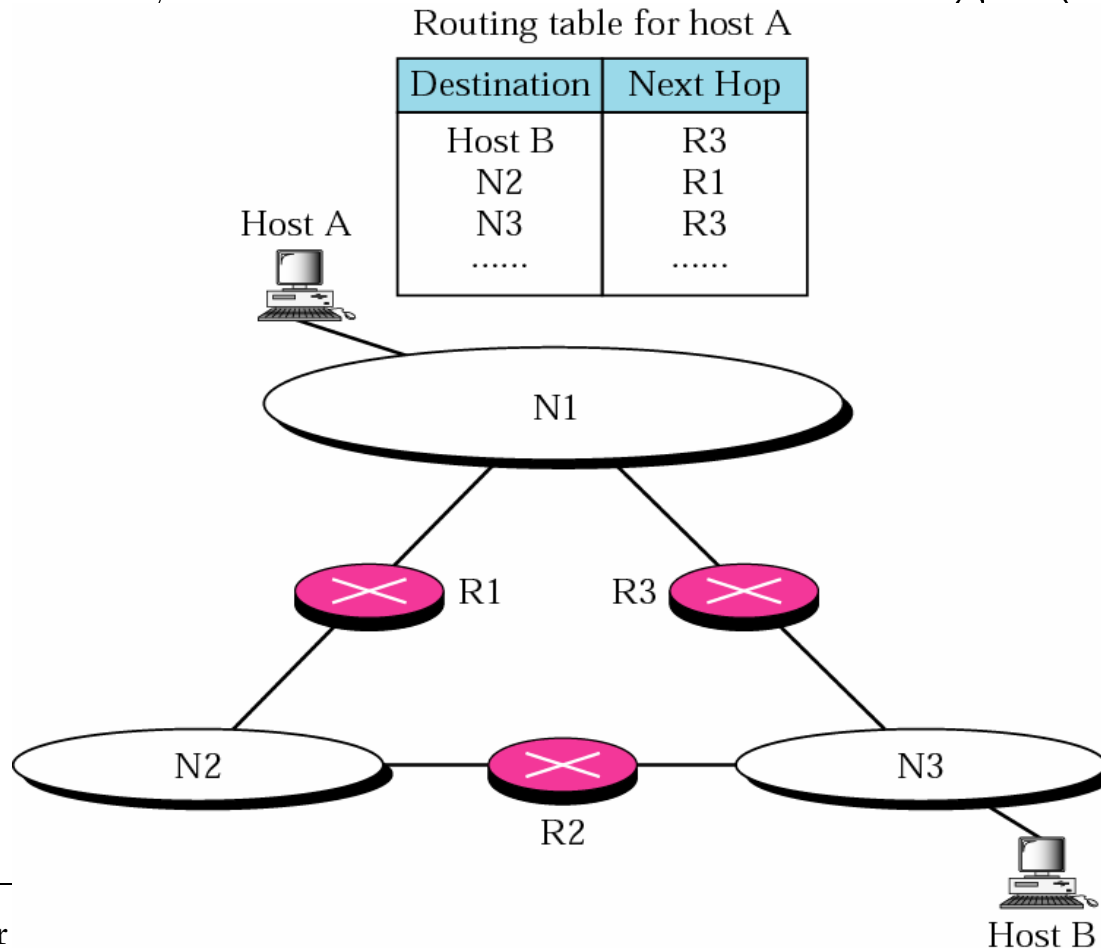
Routing table for host S based on network-specific method

| Destination | Next Hop |
|-------------|----------|
| N2 | R1 |

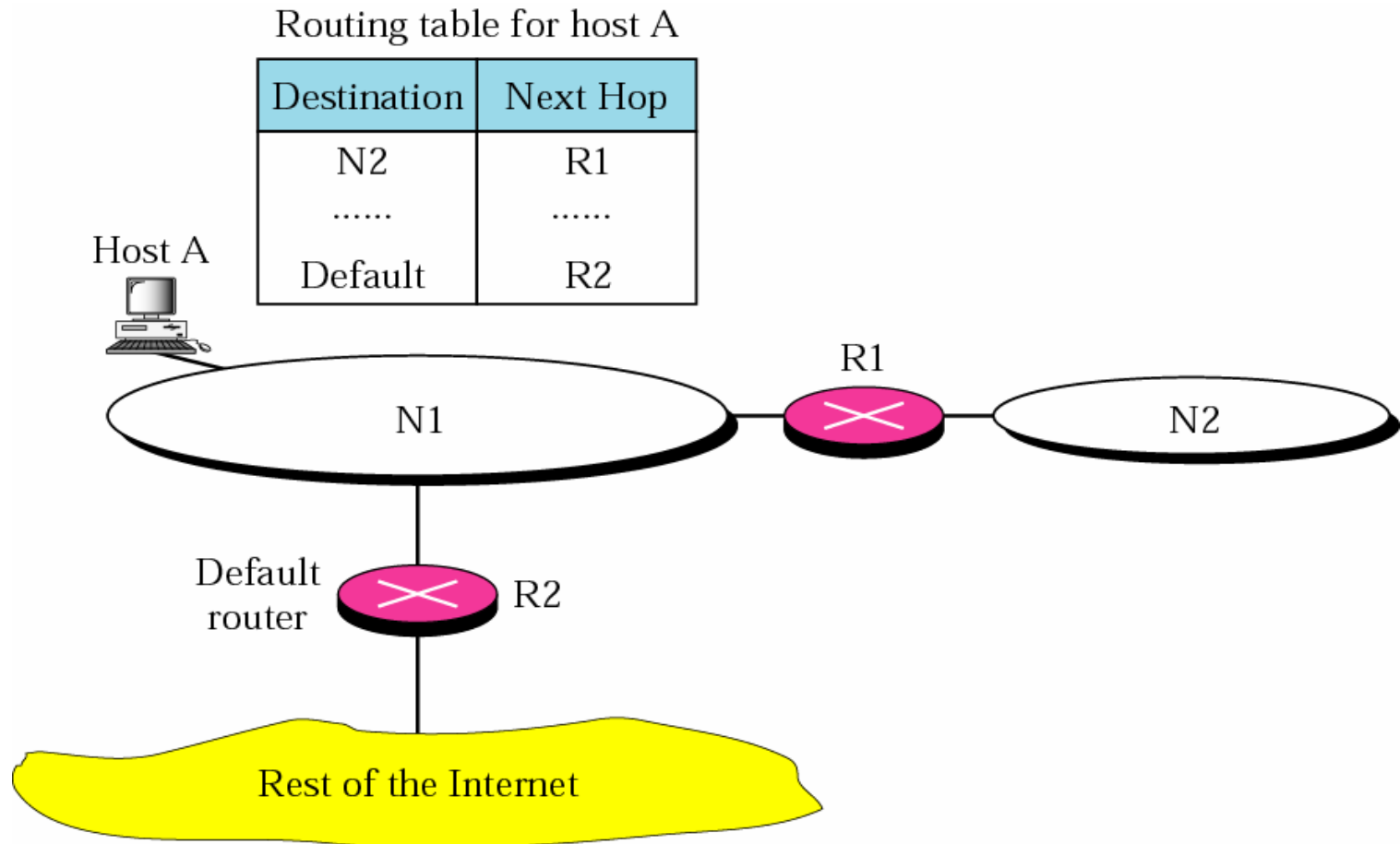


Host-Specific Method

- ▲ The destination host address is given in the routing table
- ▲ Here efficiency is sacrificed for other advantages (strict-routing)



Default Method



Forwarding with Classful Addressing

Forwarding without subnetting

- ▲ Maintain one table each for the unicast classes (A, B and C)
- ▲ If multicasting is supported, we can have a fourth table for class D

Forwarding without subnetting

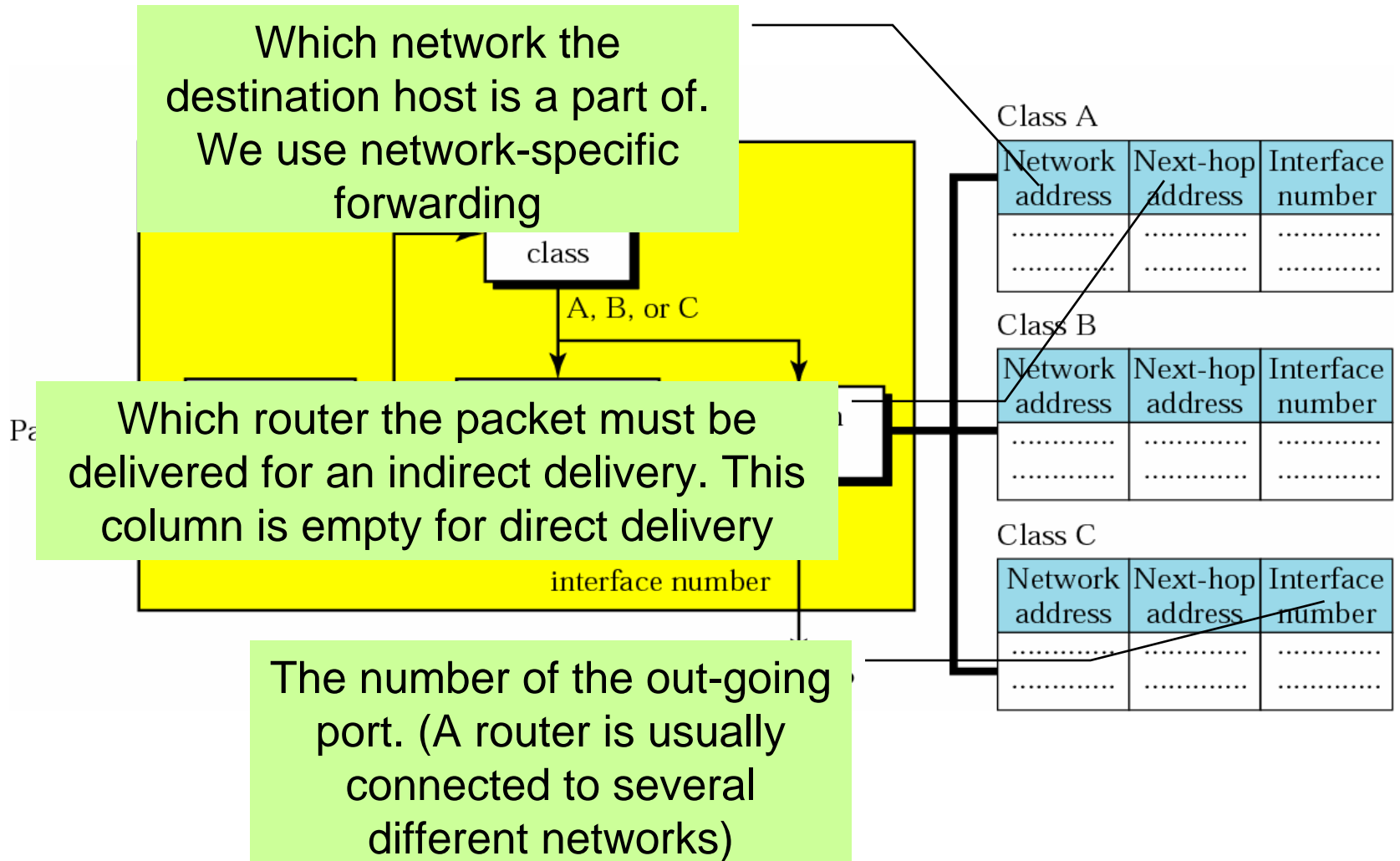
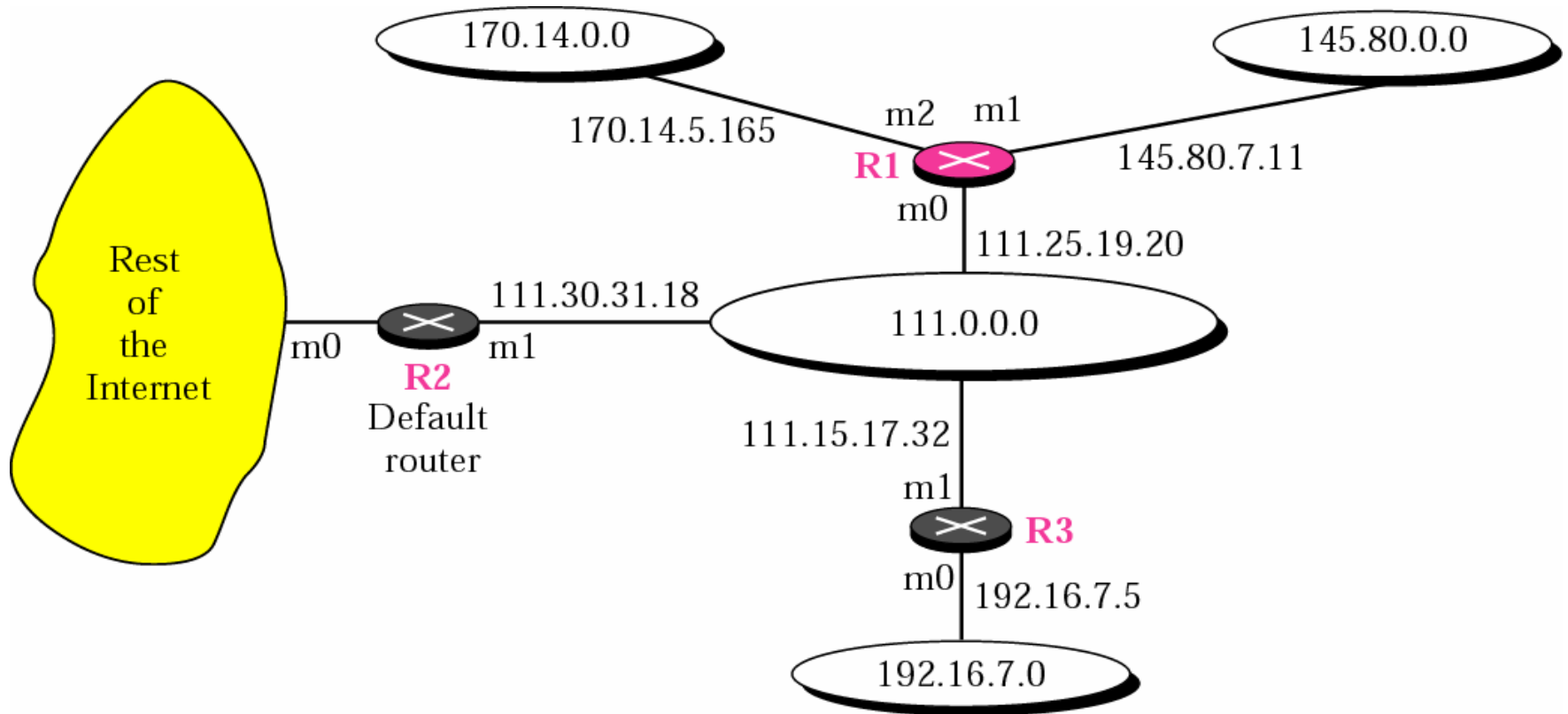


Figure shows an imaginary part of the Internet. Show the routing tables for router R1.



Class A

| Network address | Next-hop address | Interface |
|-----------------|------------------|-----------|
| 111.0.0.0 | ----- | m0 |

Class B

| Network address | Next-hop address | Interface |
|-----------------|------------------|-----------|
| 145.80.0.0 | ----- | m1 |
| 170.14.0.0 | ----- | m2 |

Class C

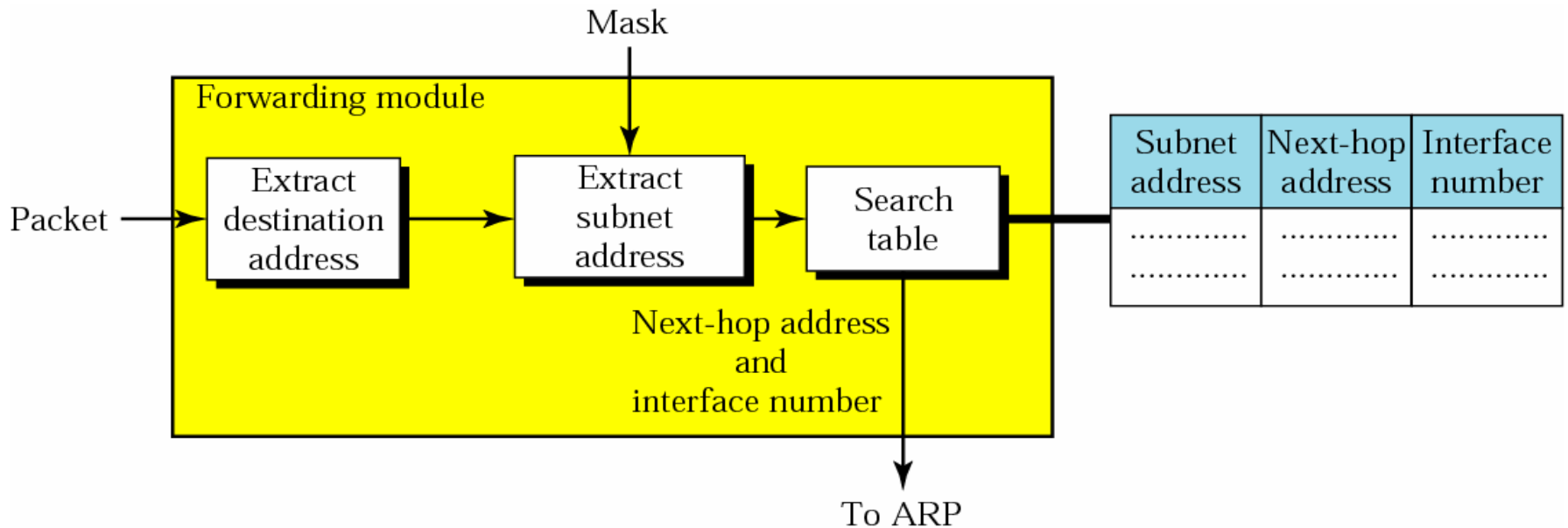
| Network address | Next-hop address | Interface |
|-----------------|------------------|-----------|
| 192.16.7.0 | 111.15.17.32 | m0 |

Default: 111.30.31.18, m0

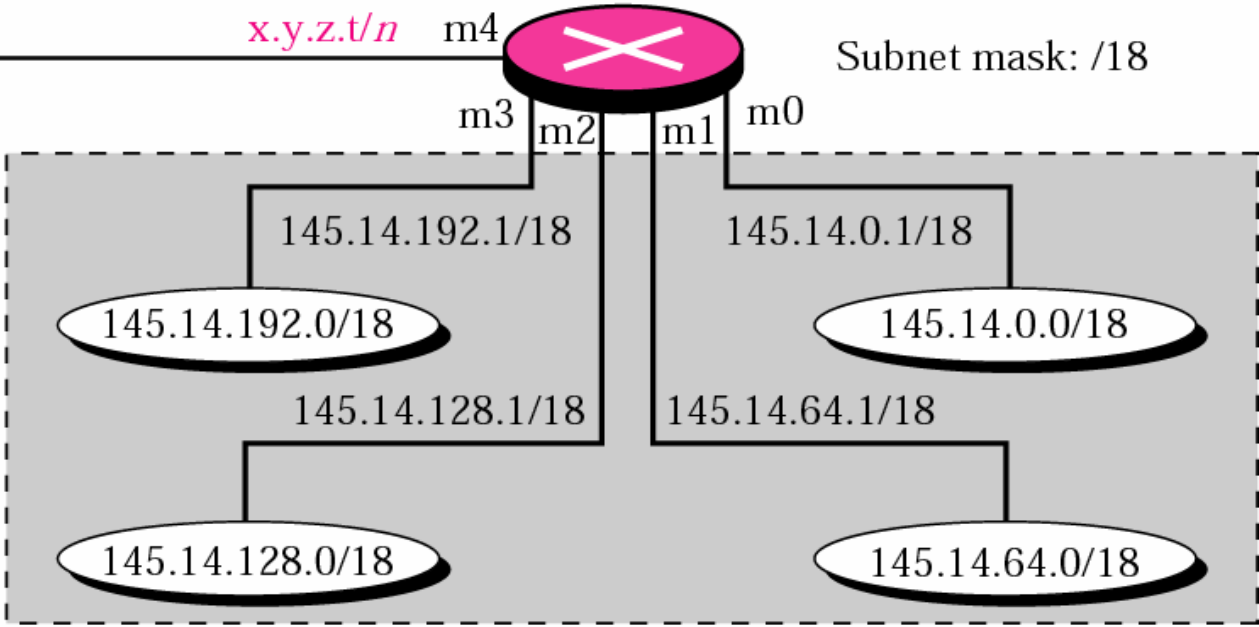
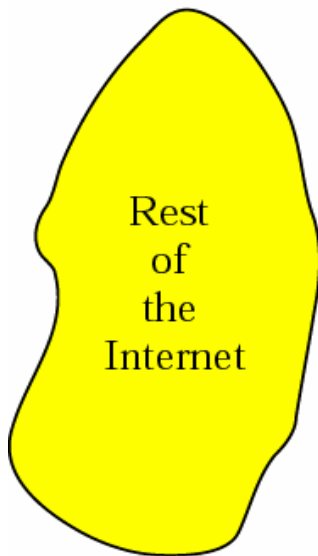
Forwarding with subnetting

- ▶ Subnetting is done inside the organization
- ▶ The routers which handle subnetting are either on the boundary or inside the boundary of the network
- ▶ If variable length subnetting is used we need several tables
- ▶ Otherwise we need only one table

Forwarding with subnetting



| Subnet address | Next-hop address | Interface number |
|----------------|------------------|------------------|
| 145.14.0.0 | ----- | m0 |
| 145.14.64.0 | ----- | m1 |
| 145.14.128.0 | ----- | m2 |
| 145.14.192.0 | ----- | m3 |
| 0.0.0.0 | default router | m4 |

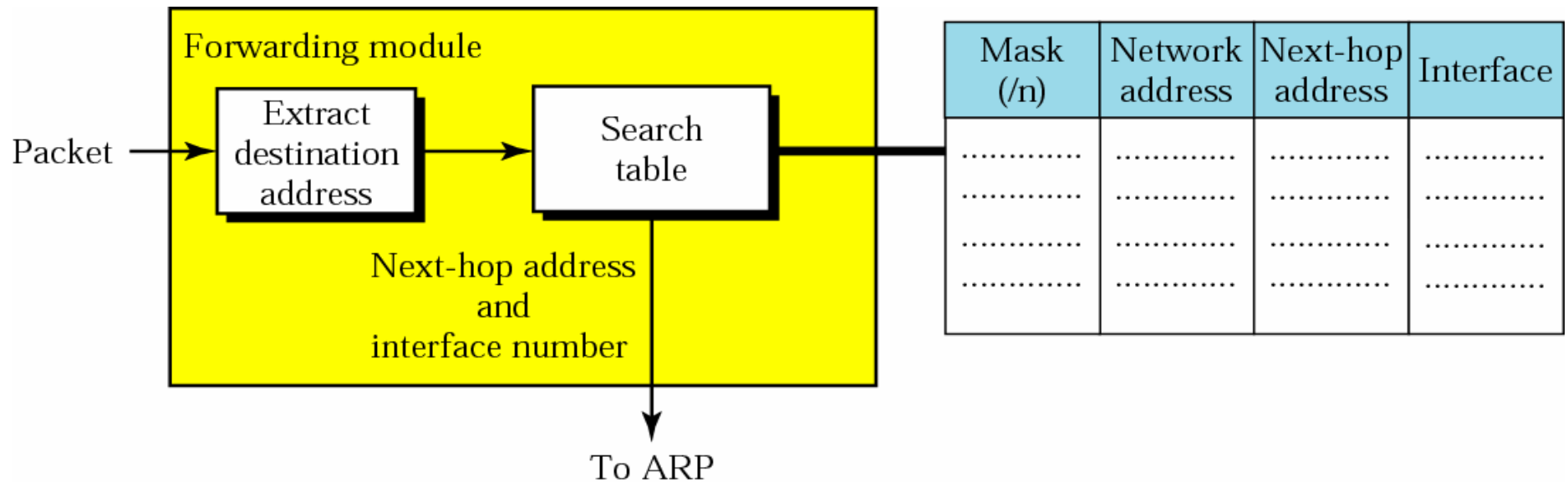


Site: 145.14.0.0/16

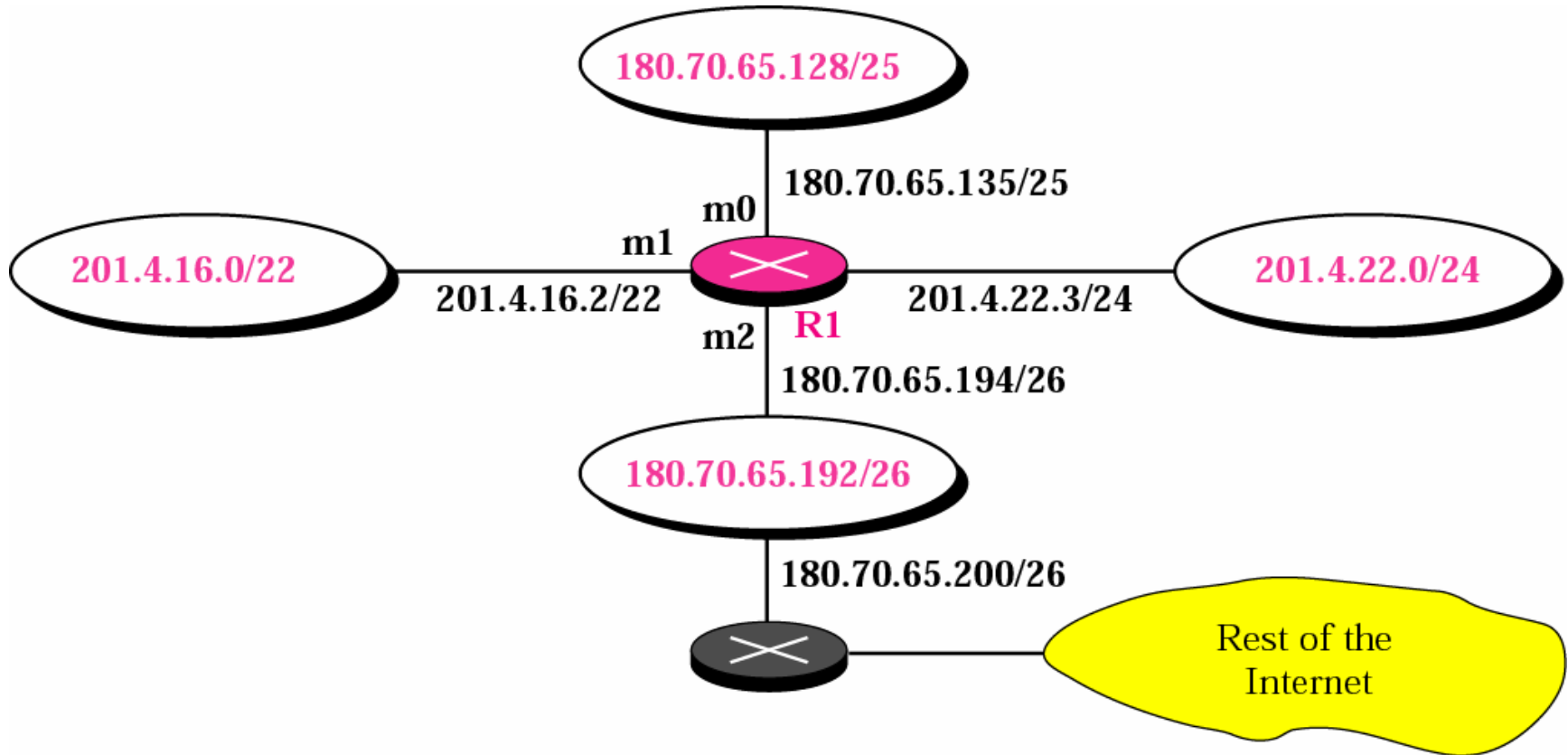
Forwarding with Classless Addressing

- ▶ Since the whole address is available for allocation, one entry for every block is required
 - As the address alone has no significance
- ▶ Information about the mask is required
- ▶ Thus we add another column for the mask

Forwarding with Classless Addressing



✦ *Make a routing table for router R1 using the configuration in Figure*



| <i>Mask</i> | <i>Network Address</i> | <i>Next Hop</i> | <i>Interface</i> |
|-------------|------------------------|-----------------|------------------|
| <i>/26</i> | 180.70.65.192 | - | m2 |
| <i>/25</i> | 180.70.65.128 | - | m0 |
| <i>/24</i> | 201.4.22.0 | - | m3 |
| <i>/22</i> | 201.4.16.0 | | m1 |
| Default | Default | 180.70.65.200 | m2 |

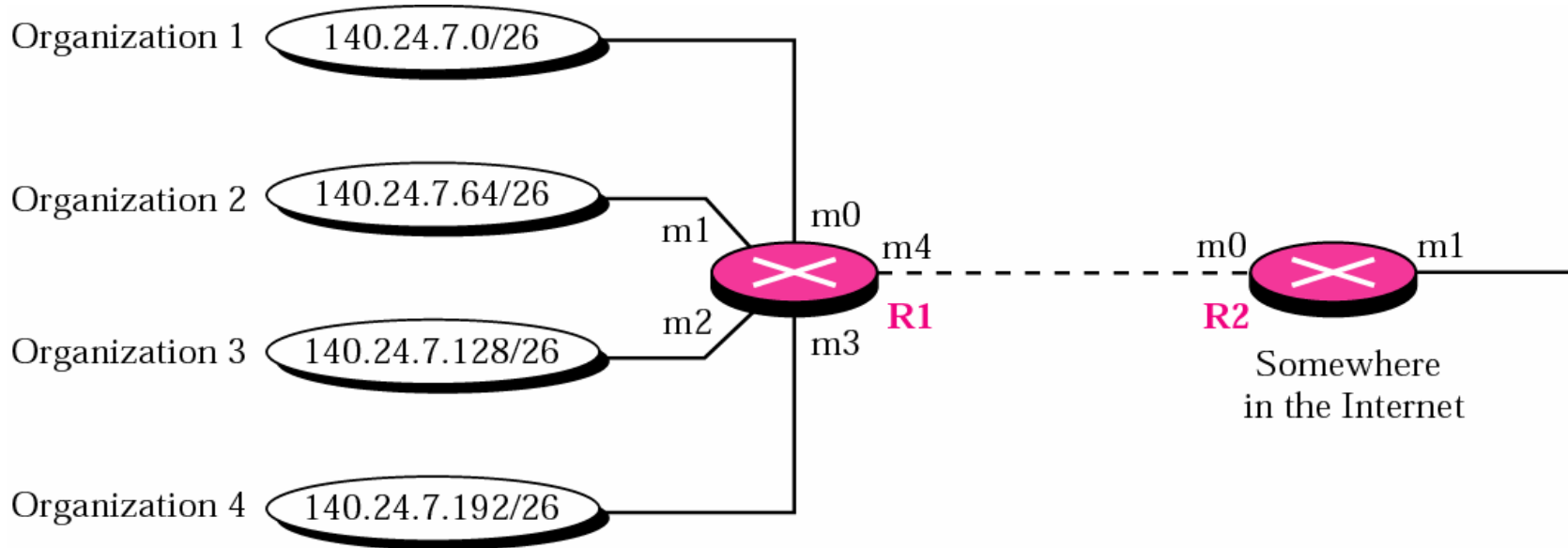
Show the forwarding process if a packet arrives at R1 in the last example with the destination address **180.70.65.140.**

- ▶ The first mask (/26) is applied to the destination address. The result is 180.70.65.128, which does not match the corresponding network address.
- ▶ The second mask (/25) is applied to the destination address. The result is 180.70.65.128, which matches the corresponding network address.

Address Aggregation

- ▲ For classful addressing there is **one** entry for every site outside the organization
- ▲ The entry defines the site (even if it is subnetted)
- ▲ With classless addressing the number of entries may increase
 - A variety of blocks
 - Similar addresses with different masks

Address Aggregation



| Mask | Network address | Next-hop address | Interface |
|------|-----------------|------------------|-----------|
| /26 | 140.24.7.0 | ----- | m0 |
| /26 | 140.24.7.64 | ----- | m1 |
| /26 | 140.24.7.128 | ----- | m2 |
| /26 | 140.24.7.192 | ----- | m3 |
| /0 | 0.0.0.0 | default router | m4 |

Routing table for R1

| Mask | Network address | Next-hop address | Interface |
|------|-----------------|------------------|-----------|
| /24 | 140.24.7.0 | ----- | m0 |
| /0 | 0.0.0.0 | default router | m1 |

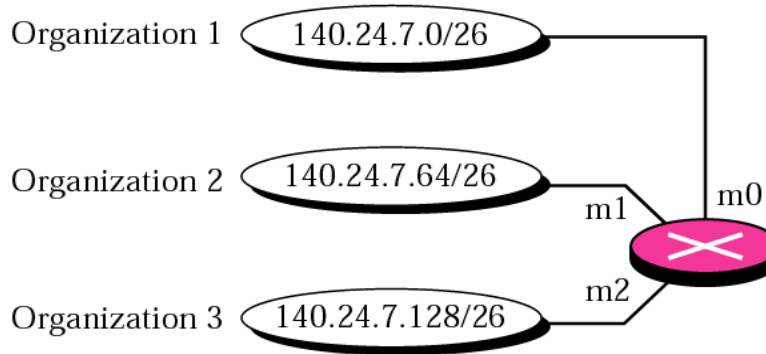
Routing table for R2

Address Aggregation

- ▶ Note that although the idea of address aggregation is similar to the idea of subnetting, **we do not have a common site here**
- ▶ The network of each organization is independent
- ▶ Also, we can have several levels of aggregation

Longest Mask Matching

- ▶ Usually address aggregation is done with sites which are close to each other –geographically
- ▶ What if an organization is not located in the same vicinity?
- ▶ Solution: longest mask matching
 - Sorted with the longest mask first

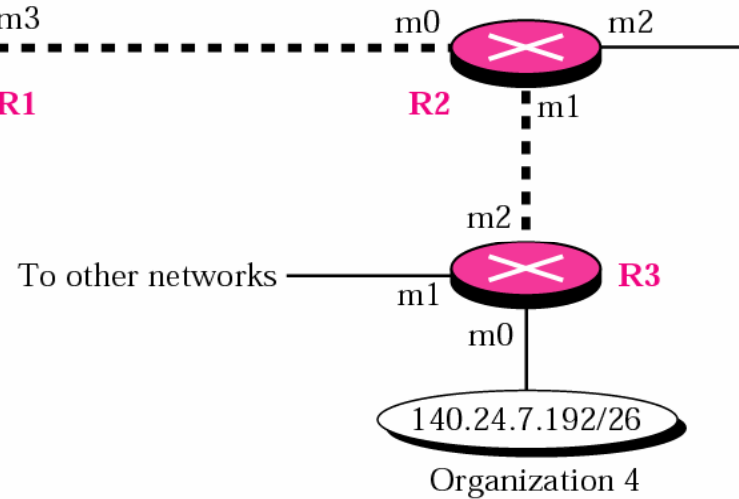


| Mask | Network address | Next-hop address | Interface |
|------|-----------------|------------------|-----------|
| /26 | 140.24.7.0 | ----- | m0 |
| /26 | 140.24.7.64 | ----- | m1 |
| /26 | 140.24.7.128 | ----- | m2 |
| /0 | 0.0.0.0 | default router | m3 |

Routing table for R1

| Mask | Network address | Next-hop address | Interface |
|------|-----------------|------------------|-----------|
| /26 | 140.24.7.192 | ----- | m1 |
| /24 | 140.24.7.0 | ----- | m0 |
| /?? | ???????? | ?????????? | m1 |
| /0 | 0.0.0.0 | default router | m2 |

Routing table for R2

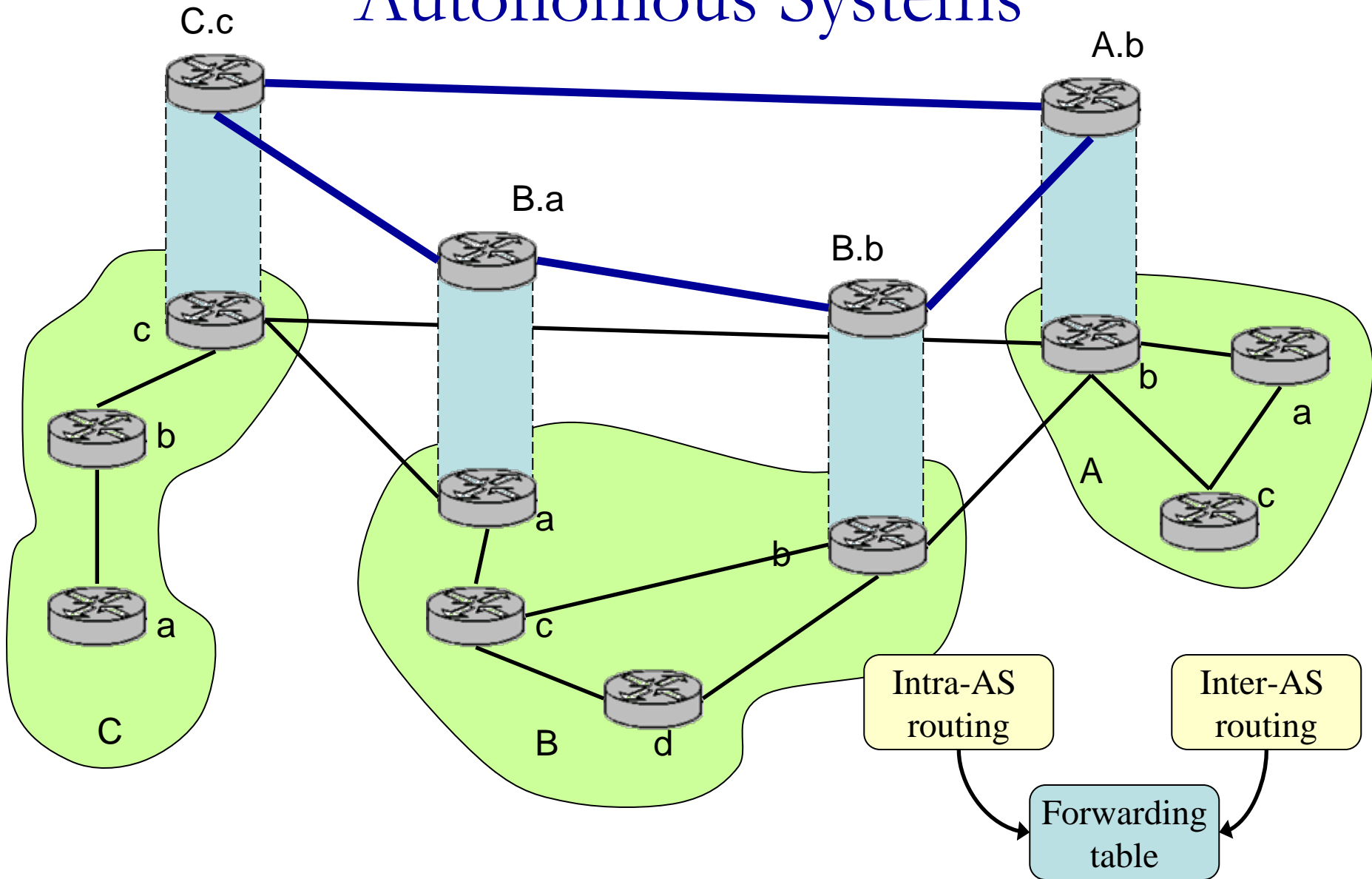


| Mask | Network address | Next-hop address | Interface |
|------|-----------------|------------------|-----------|
| /26 | 140.24.7.192 | ----- | m0 |
| /?? | ???????? | ?????????? | m1 |
| /0 | 0.0.0.0 | default router | m2 |

Routing table for R3

Autonomous Systems & Hierarchical Routing

Autonomous Systems



Questions?

That's all folks!